# Introduction to Information Retrieval

**Crawling and Duplicates** 

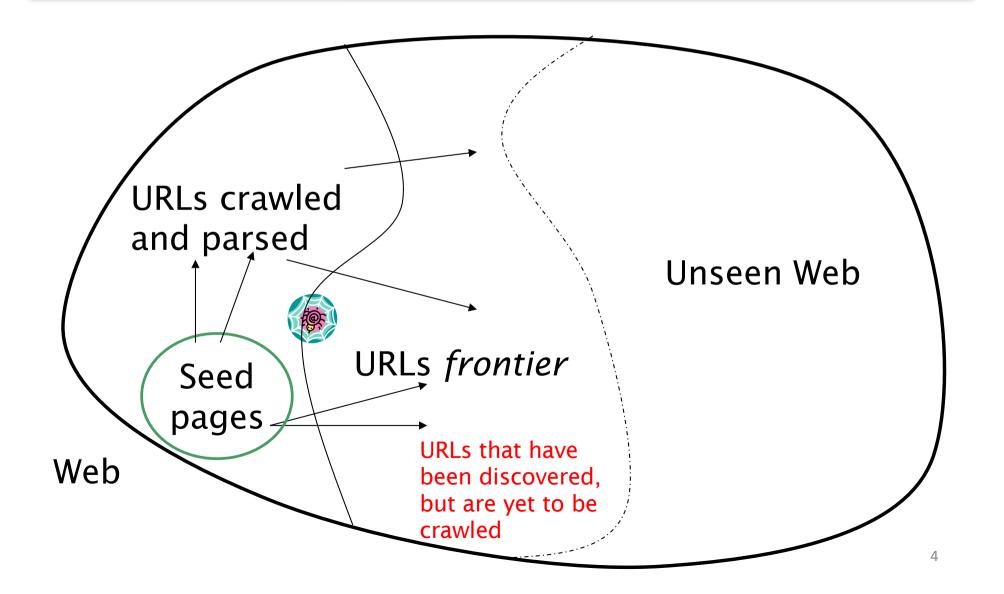
# This lecture

- Web Crawling
- (Near) duplicate detection

# **Basic crawler operation**

- Begin with known "seed" URLs
- Fetch and parse them
  - Extract URLs they point to
  - Place the extracted URLs on a queue
- Fetch each URL on the queue and repeat
- Breadth First crawling

## Crawling picture



Sec. 20.2

## Simple picture – complications

- Web crawling isn't feasible with one machine
  - All of the above steps are usually distributed
- Malicious pages
  - Spam pages
  - Spider traps
- Even non-malicious pages pose challenges
  - Latency/bandwidth to remote servers vary
  - Webmasters' stipulations
    - How "deep" should you crawl a site's URL hierarchy?
  - Site mirrors and duplicate pages
- Politeness don't hit a server too often

#### What any crawler *must* do

Be <u>Polite</u>: Respect implicit and explicit politeness considerations

Be <u>Robust</u>: Be immune to spider traps and other malicious behavior from web servers

Sec. 20.2

## Explicit and implicit politeness

- <u>Explicit politeness</u>: specifications from webmasters on what portions of a site can be crawled
  - robots.txt (see next slide)
- <u>Implicit politeness</u>: even with no specification, avoid hitting any site too often

#### Robots.txt

- Protocol for giving spiders ("robots") limited access to a website, originally from 1994
- Website announces its request on what can(not) be crawled
  - For a server, create a file / robots.txt
  - This file specifies access restrictions

Details: <u>www.robotstxt.org/robotstxt.html</u>

## What any crawler *should* do

- Be capable of <u>distributed</u> operation: designed to run on multiple distributed machines
- Be <u>scalable</u>: designed to increase the crawl rate by adding more machines

 <u>Performance/efficiency</u>: permit full use of available processing and network resources

## What any crawler should do

- Fetch pages of "higher <u>quality</u>" first
- <u>Continuous</u> operation: Continue fetching fresh copies of a previously fetched page
- Extensible: Adapt to new data formats, protocols

## URL frontier

- URLs that have been discovered, but are yet to be crawled
- Can include multiple pages from the same host
- Must avoid trying to fetch them all at the same time
- Must try to keep all crawling threads busy

### Processing steps in crawling

- Pick a URL from the frontier
- Fetch the document at the URL
- Parse the URL
  - Extract links from it to other docs (URLs)
- Check if URL has content already seen
  - If not, add to indexes
- For each extracted URL
  - Ensure it passes certain URL filter tests
  - Check if it is already in the frontier (duplicate URL elimination)

E.g., only crawl .edu, obey robots.txt, etc.

## Parsing: URL normalization

- When a fetched document is parsed, some of the extracted links are *relative* URLs
- E.g., <u>http://en.wikipedia.org/wiki/Main\_Page</u> has a relative link to /wiki/Wikipedia:General\_disclaimer which is the same as the absolute URL <a href="http://en.wikipedia.org/wiki/Wikipedia:General\_disclaimer">http://en.wikipedia.org/wiki/Wikipedia:General\_disclaimer</a>
- During parsing, must normalize (expand) such relative URLs

## Content seen?

- Duplication is widespread on the web
- If the page just fetched is already in the index, do not further process it
- This is verified using document fingerprints or <u>shingles</u>
  - Second part of this lecture

## Distributing the crawler

- Run multiple crawl threads, under different processes – potentially at different nodes
  - May be geographically distributed nodes
- Partition hosts being crawled into nodes

#### URL frontier: two main considerations

- Politeness: do not hit a web server too frequently
- <u>Freshness</u>: crawl some pages more often than others
  - E.g., pages (such as News sites) whose content changes often

These goals may conflict with each other.

(E.g., simple priority queue fails – many links out of a page go to its own site, creating a burst of accesses to that site.)

## Politeness – challenges

- Even if we restrict only one thread to fetch from a host, can hit it repeatedly
- Common heuristic: insert time gap between successive requests to a host that is >> time for most recent fetch from that host

# Introduction to Information Retrieval

Near duplicate document detection

## Duplicate documents

- The web is full of duplicated content
- Strict duplicate detection = exact match
  - Not as common
- But many, many cases of near duplicates
  - E.g., Last modified date the only difference between two copies of a page

#### Duplicate/Near-Duplicate Detection

- Duplication: Exact match can be detected with fingerprints
- Near-Duplication: Approximate match
  - Overview
    - Compute syntactic similarity with an edit-distance measure
    - Use similarity threshold to detect near-duplicates, e.g., Similarity > 80% => Documents are "near duplicates"

## **Computing Similarity**

- Features:
  - Segments of a document (natural or artificial breakpoints)
  - Shingles (Word N-Grams)
  - a rose is a rose is a rose → 4-grams are

```
a_rose_is_a
rose_is_a_rose
is_a_rose_is
```

- Similarity Measure between two docs (= sets of shingles)
  - Jaccard cooefficient: (Size\_of\_Intersection / Size\_of\_Union)

## Shingles + Set Intersection

- Computing <u>exact</u> set intersection of shingles between <u>all</u> pairs of documents is expensive
- Approximate using a cleverly chosen subset of shingles from each (a *sketch*)
- Estimate (size\_of\_intersection / size\_of\_union) based on a short sketch

$$\begin{array}{c} Doc \\ A \end{array} \rightarrow Shingle set A \rightarrow Sketch A \\ \hline Doc \\ B \end{array} \rightarrow Shingle set B \rightarrow Sketch B \end{array}$$
 Jaccard

## Sketch of a document

- Create a "sketch vector" (of size ~200) for each document
  - Documents that share ≥ t (say 80%) corresponding vector elements are deemed near duplicates
  - For doc D, sketch<sub>D</sub>[i] is as follows:
    - Let f map all shingles in the universe to 1..2<sup>m</sup> (e.g., f = fingerprinting)
    - Let  $\pi_i$  be a *random permutation* on  $1..2^m$
    - Pick MIN { $\pi_i(f(s))$ } over all shingles *s* in *D*

See details in book

## **Final notes**

- Shingling is a randomized algorithm
  - It will give us the right (wrong) answer with some probability on *any input*
- We've described how to detect near duplication in a pair of documents
- In "real life" we'll have to concurrently look at many pairs
  - See text book for details