Introduction to Information Retrieval

Lectures 4: Skip Pointers, Phrase Queries, Positional Indexing

Introduction to Information Retrieval

Faster postings merges: Skip pointers/Skip lists

Recall basic merge

 Walk through the two postings simultaneously, in time linear in the total number of postings entries

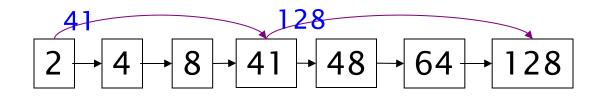
$$2 \rightarrow 8 \qquad \qquad 2 \rightarrow 4 \rightarrow 8 \rightarrow 41 \rightarrow 48 \rightarrow 64 \rightarrow 128 \quad Brutus$$
$$1 \rightarrow 2 \rightarrow 3 \rightarrow 8 \rightarrow 11 \rightarrow 17 \rightarrow 21 \rightarrow 31 \quad Caesar$$

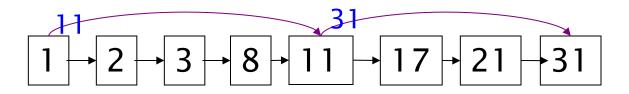
If the list lengths are *m* and *n*, the merge takes O(*m*+*n*) operations.

Can we do better? Yes (if the index isn't changing too fast). Introduction to Information Retrieval

Augment postings with skip pointers (at indexing time)

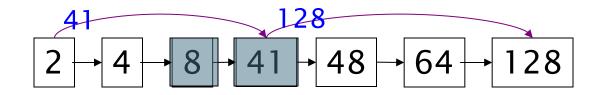
Sec. 2.3

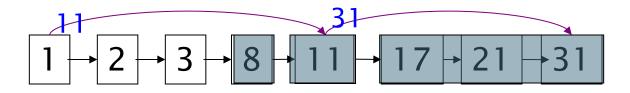




- Why?
- <u>To skip postings that will not figure in the search</u> <u>results.</u>
- How?
- Where do we place skip pointers?

Query processing with skip pointers





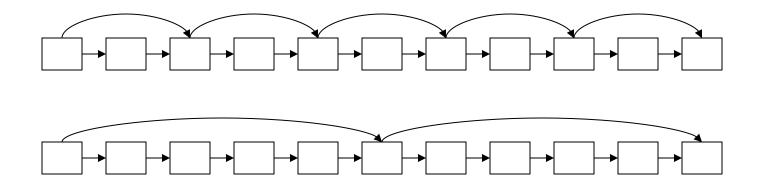
Suppose we've stepped through the lists until we process 8 on each list. We match it and advance.

We then have **41** and **11** on the lower. **11** is smaller.

But the skip successor of 11 on the lower list is 31, so we can skip ahead past the intervening postings.

Where do we place skips?

- Tradeoff:
 - More skips → shorter skip spans ⇒ more likely to skip.
 But lots of comparisons to skip pointers.
 - Fewer skips → few pointer comparison, but then long skip spans ⇒ few successful skips



Placing skips

- Simple heuristic: for postings of length *L*, use \sqrt{L} evenly-spaced skip pointers [Moffat and Zobel 1996]
- Easy if the index is relatively static; harder if *L* keeps changing because of updates.
- This definitely used to help; with modern hardware it may not unless you're memory-based [Bahle et al. 2002]
 - The I/O cost of loading a bigger postings list can outweigh the gains from quicker in memory merging!

Introduction to Information Retrieval

Handling phrase queries

Phrase queries

We want to answer a query such as [stanford university] – as a phrase.

Thus The inventor Stanford Ovshinsky never went to university should not be a match.

The concept of phrase query has proven easily understood by users.

About 10% of web queries are phrase queries.

Consequence for inverted index: it no longer suffices to store docIDs in postings lists for terms.

- Two ways of extending the inverted index:
 - biword index
 - positional index

Biword indexes

- Index every consecutive pair of terms in the text as a phrase.
- •For example, *Friends, Romans, Countrymen* would generate two biwords: *"friends romans"* and *"romans countrymen"*
- Each of these biwords is now a vocabulary term.
- Two-word phrases can now easily be answered.

Longer phrase queries

A long phrase like "stanford university palo alto" can be represented as the Boolean query "STANFORD UNIVERSITY" AND "UNIVERSITY PALO" AND "PALO ALTO"

Does this always guarantee the correct match? -- We need to do post-filtering of hits to identify subset that actually contains the 4-word phrase.

What about phrases like, "abolition of slavery"?

Extended biwords

- Parse each document and perform part-of-speech tagging
- Bucket the terms into (say) nouns (N) and articles/prepositions
 (X)
- Now deem any string of terms of the form NX*N to be an extended biword
- Examples: catcher in the rye

Ν

king of Denmark

N X N

Include extended biwords in the term vocabulary

X X N

Queries are processed accordingly

Issues with biword indexes

Why are biword indexes rarely used?

- False positives, as noted above
- Index blowup due to very large term vocabulary

•What can be an alternative?

Positional indexes

 Positional indexes are a more efficient alternative to biword indexes.

Postings lists in a nonpositional index: each posting is just a docID

Postings lists in a positional index: each posting is a docID and a list of positions

Positional indexes: Example

Query: "to₁ be₂ or₃ not₄ to₅ be₆" то, 993427: < 1: <7, 18, 33, 72, 86, 231>; **2**: <1, 17, 74, 222, 255>; 4: <8, 16, 190, 429, 433>; 5: <363, 367>; 7: <13, 23, 191>; . . . > BE, 178239: < 1: <17, 25>; 4: <17, 191, 291, 430, 434>; 5: <14, 19, 101>; . . . >

Proximity search

- •We just saw how to use a positional index for phrase searches.
- Can we also use it for proximity search?
- For example: employment /4 place
- Find all documents that contain EMPLOYMENT and PLACE within 4 words of each other.
 - Employment agencies that place healthcare workers are seeing growth is a hit.
 - Employment agencies that have learned to adapt now place healthcare workers is not a hit.

Proximity search

- Use the positional index
- Simplest algorithm: look at cross-product of positions of (i) EMPLOYMENT in document and (ii) PLACE in document
- Very inefficient for frequent words, especially stop words
- Note that we want to return the actual matching positions, not just a list of documents.

Combination scheme

- Biword indexes and positional indexes can be profitably combined.
- Many biwords are extremely frequent: Michael Jackson etc
- For these biwords, increased speed compared to positional postings intersection is substantial.
- Combination scheme: Include frequent biwords as vocabulary terms in the index. Do all other phrases by positional intersection.
- Williams et al. (2004) evaluate a more sophisticated mixed indexing scheme – Next Word Index. Faster than a positional index, at a cost of 26% more space for index.